University of Ljubljana, Faculty of Computer and Information Science

Intelligent agents





Prof Dr Marko Robnik-Šikonja, Intelligent Systems, January 2020

Contents

- types of agents and agent architectures
- multiagent learning
- # distributed constraint satisfaction
- # distributed shortest path finding

Literature

Yoav Shoham, Kevin Leyton-Brown: <u>Multiagent</u> <u>Systems: Algorithmic, Game-Theoretic, and Logical</u> <u>Foundations</u>. Cambridge University Press, 2009

This lecture is only a teaser, we cover part of Chapters 1 and 2, but things get useful with game theory.

Some terminology

- agent: many definitions, many areas
- Intelligent agent: autonomous entity which observes and acts upon an environment and directs its activity towards achieving goals
- Biological, artificial
- Software agents, intelligent agents
- Types and features
- Agent architectures
- A fundamental concept in intelligent systems

Agent

- Control system
- Sensors
- Actuators
- Environment

- Autonomous: acts
 independently and makes
 its own decisions
- Social ability: can interact with other agents
- Reactivity: reacts to stimuli
- Proactivity: pursues its own goals and acts in its own (self-)interest

Multiagent systems

Distributed program solving

- x autonomous,
- % flexible,
- **☆** collectively organized actions
- \approx (goal oriented).
- Examples: drones, kilobots, helicopters, boats etc, see any of the videos on this topic: <u>an example</u>

Interactive environment

- agent percepts information
- actions affect the environment
- internet, game playing, robotics (e.g., robosoccer)

Autonomous and flexible

- Without explicit directives
- Control
- Model of the environment
- (learning from experience)
- Responsiveness,
- Planning, goal oriented



Collective actions

- Interactions with agents and humans
- Goal oriented cooperation: strategies, negotiation, specialization
- Distributed, asynchronous
- Objects and agents



Applications

- Production
- Automatic process control
- Telecommunications
- Information management
- e-business
- Interactive games
- Services





Types of agents

- reactive agents
- collaborative agents
- interface agents
- mobile agents
- information-gathering agents
- hybrid agents

Reactive agents

- Response according to pre-specified rules
- Mail sorter, spam filter, calendar management
- Learning and revision of rules

Goal oriented agents

- Following the goal
- Planning and search
- Tickets, products

Utility based agents

- Utility functions
- Goals and utility maximization, Multiobjective decision making
- Rationality (e.g., in games one can loose on purpose)
- Kahneman, Tversky: (A. 100% 3000, B. 80% 4000
 C. 20% 4000 D. 25% 3000),

Interface agents

- Personal assistants,
- Learning
- Tutoring systems, preference learning in search

Mobile agents

Physical and virtual mobility

virus, supervision program

Information-gathering agents

- Internet, intranet, mail
- Precision, recall
- Learning
- Noisy data, data relevance

Collaborative agents

Weakly interactive agents (ants, genetic algorithms)

redundancy, parallelism

Agent architecture

 a blueprint for software agents and intelligent control systems, depicting the arrangement of components

Subsumption architecture

- Brooks, 1985,
- intelligence without representation
- Multilevel
- Each level follows its own goal
- Higher levels can block lower levels



- Each level contains its own rules, e.g., if-then rules
- Adding new levels is easy
- Debugging is hard

BDI architecture

- BDI (Belief Desire Intention)
- # planning
- Bold and cautious

Blackboard architecture

Sharing common work area

- Specialization
- Coordination
- Threads



Mobile architecture

- Agent-based modeling software (many frameworks)
- Early example:
 Aglets, IBM 1990



- JAVA Agent DEvelopment Framework
- * an open source platform for peer-to-peer agent based applications
- conforms to FIPA standard (Foundation for Intelligent Physical Agents)

Horizontal and vertical architectures



Environment

- Deterministic
- Nondeterministic

Learning agents

Learning is an adaptation

Multiagent learning

- ℅ Centralized
- ☆ Distributed

Robotic agents

- Complexity of real-world environment
- Risk management
- Industrial robots
- Robot explorers (Mars, autonomy, moving around, insects)

Robotics : Intelligent Systems view

Testbed for intelligent systems





fixed (industrial, robotic hand with several degrees of freedom)



legs (1,2,4,6,8)



Taxonomy



Underwater and amphibious (fish, crabs, worms)



Taxonomy

Airborne (drones, quadropters, satellites)





Polymorphic, swarmsPhysical and softbots

Sensors

- vision (electromagnetic waves)
- hearing (air)
- Taste and smell (chemical receptors)
- touch (pressure)
- echolocation (ultrasound)
- electroception (electric stimuli, current and field)
- magnetoception (detect magnetism, magnetic field)
- equilibrioception (balance, acceleration)
- thermoception (temperature)





Wheels, legs, motors,

Hands, grasps,

Control system

- * reactive
- Subsumption
- Neural networks, evolutionary approaches




Planning

- Essential component of intelligent behavior
- Anytime planning: always ready, but improves with time
- Cell decomposition
- Potential field
- Landmarks
- Visibility graph



Tools

- Robotic languages
- Simulators
- Robot operating systems (ROS)

Distributed computing with agents

- Agents collaborate to achieve a common goal defined by central authority
- Autonomous agents, only local communication
- The goal is to find a solution with global constraints
- The task: prepare an algorithm for the agents

An example: sensor network

- Limited
 computational
 resources
- Local communication
- Global
 constraints and
 solutions



Constraint satisfaction problem

- Set of variables with their domains and constraints on values taken by the variables
- The task: assign values to the variables satisfying all constraints or proclaim that there is no such an assignment
- Several applications: planning, vision, NLP, theorem proving, scheduling

An example: find non-overlapping frequencies for the sensors

- Three sensors
- Overlapping reach
- The task: assign nonoverlapping frequencies to the sensors from the domain of allowed frequencies



An example



- Equivalent to graph coloring
- * Set of variables $X = \{X_1, X_2, X_3\}$
- Domain D_i for each variables is {red, blue, green}
- * Set of constraints $\{X_1 \neq X_2, X_1 \neq X_3, X_3 \neq X_2\}$

Constraint satisfaction terminology

- Variable assignment
 - Legal, illegal

Solution

 Distributed constraint satisfaction: each agent is a variable, the solution is to be found without central control

Domain pruning algorithms

- Nodes communicate with neighbors to prune forbidden values from their domains
- arc consistency algorithm
- Each vertex X_i with domain D_i repeatedly executes the program for each of its neighbors X_i
- void revise(x_i, x_j) {
 - foreach ($v_i \in D_i$) if (there is no value $v_j \in D_j$ such that v_i is consistent with v_j) $D_i = D_i - \{v_i\}$

Arc consistency

- Stop when one of domains is empty (no solution), or no more eliminations takes place.
- If there is a single value left in each domain, we have a solution; otherwise the result is inconclusive: we do not know if the solution exists.
- Algorithm terminates and is sound (the solution if found, is correct), but it is not complete (no guarantee that the solution will be found).

An example of domain pruning a)

- First only messages to node X1 are efficient, therefore X2 and X3 eliminate value red
- X2={blue} X3={blue, green}
- Next X₃ can eliminate blue;
- The result is correct



An example of domain pruning b)

- As before X₂ and X₃
 eliminate *red*
- X2={blue} X3={blue}
- Next both X2 and X₃ eliminate *blue;*
- Empty domain is left, so they proclaim there is no solution.



An example of domain pruning c), d)

No node can eliminate any value

Inconclusive termination





Equivalence to logic resolution

- Arc consistency is too weak, can be used as preprocessing
- Value elimination is equivalent to unit resolution in propositional logic
- Inference rule
 A1 $\neg(A_1 \land A_2 \land \dots \land A_n)$

 $\neg (A_2 \land ... \land A_n)$

 We write constraints as forbidden value combinations called Nogoods, e.g., x_1 =red $\land x_2$ =red

 $x_1 = red$ $\neg(x_1 = red \land x_2 = red)$

 $\neg(x_2 = red)$

Hyper-resolution

* A generalization of unit resolution $A_1 \lor A_2 \lor \cdots \lor A_m$ $\neg (A_1 \land A_{1,1} \land A_{1,2} \land \cdots)$ $\neg (A_2 \land A_{2,1} \land A_{2,2} \land \cdots)$... $\neg (A_m \land A_{m,1} \land A_{m,2} \land \cdots)$

$$\neg (A_{1,1} \land \cdots \land A_{2,1} \land \cdots \land A_{m,1} \land \cdots)$$

- Sound and complete for propositional logic
- at least one of the literals in the top disjunction A₁ V A₂ V · · · V A_m is true, therefore the conjunction of all the remaining literals in the negated terms has to fail, too

Hyper-resolution algorithm

- each agent repeatedly generates new constraints for his neighbors, notifies them of these new constraints, and prunes his own domain based on new constraints passed to him by his neighbors.
- * NG_i is the set of all Nogoods of which agent i is aware
- * NG*_j is a set of new Nogoods communicated from agent j to agent i.

```
void reviseHR(NG<sub>i</sub>, NG<sup>*</sup><sub>i</sub>) {
    do {
       NG_{i} \leftarrow NG_{i} \cup NG_{j}^{*}
NG_{i}^{*} \leftarrow hyperresolution(NG_{i}, D_{i})
       if (NG_{i}^{*} \neq \{\})
         NG_i \leftarrow NG_i \cup NG^*_i
         send the Nogoods NG<sup>*</sup>, to all neighbours of i
         if ( \{\} \in NG^* \}
            stop
    } while (there is a change in NG<sub>i</sub>)
}
```

algoritem terminates after finite number of steps

If the solution exists, the algorithms finds it

Hyper-resolution for c)

 X₁ has initially the following constraints in its Nogoods: {x₁ = red, x₂ = red}, {x₁ = red, x₃ = red}, {x₁ = blue, x₂ = blue}, {x₁ = blue, x₃ = blue}



- X_1 can be assigned values $x_1 = \text{red } V x_1 = \text{blue}$.
- With hyper-resolution X₁ can reason

 $x_{1} = red \lor x_{1} = blue$ $\neg(x_{1} = red \land x_{2} = red)$ $\neg(x_{1} = blue \land x_{3} = blue)$

 $\neg(x_2 = red \land x_3 = blue)$

And adds constraints {x2 = red, x3 = blue} to its Nogoods

Hyper-resolution for c)

- Similarly it adds
 {x₂ = blue, x₃ = red} to its Nogoods
- x₁ sends both Nogoods to its neighbors x₂ and x₃
- * x_2 can reason (based on its domain, Nogoods and received inferences) $x_2 = red \lor x_2 = blue$ $\neg(x_2 = red \land x_3 = blue)$ $\neg(x_2 = blue \land x_3 = blue)$

 $\neg(x_3 = blue)$

- ***** Based on the other received Nogood x_2 constructs $\neg(x_3 = red)$
- When both Nogoods are send to the neighbor x₃, x₃ generates {} and algorithms terminates proclaiming that no solution exists.



Weaknesses of hyper-resolution

- Number of generated Nogoods can be very large
- Asynchronous and parallel processing would generate even more Nogoods.
- The problem lies in the least-commitment nature of these algorithms; they are restricted to removing only provably impossible value combinations.
- The alternative is to explore a subset of the space, making tentative value selections for variables, and backtracking when necessary.

Heuristic search with constraints

- Centralized trial and error
- Sort variables, e.g., x₁, x₂, ..., x_n
- * Call chooseValue(x_1 , {}), with values { v_1 , v_2 , ..., v_{i-1} } already assigned to { x_1 , x_2 , ..., x_{i-1} }

procedure ChooseValue $(x_i, \{v_1, v_2, \dots, v_{i-1}\})$ $v_i \leftarrow$ value from the domain of x_i that is consistent with $\{v_1, v_2, v_{i-1}\}$ **if** no such value exists **then**

- backtrack¹
- else if i = n then
 - stop
- else

ChooseValue($x_{i+1}, \{v_1, v_2, ..., v_i\}$)

Weaknesses of chronological backtracking

- Exhaustive search
- Agents work sequentially

Naïve parallel asynchronous solution

executed by all agents in parallel and asynchronously

select a value from your domain

repeat

if your current value is consistent with the current values of your neighbors, or if none of the values in your domain are consistent with them **then**

do nothing

else

select a value in your domain that is consistent with those of your neighbors and notify your neighbors of your new value

until there is no change in your value

Correct but incomplete solution: it may not terminate, it may not find a solution

Asynchronous backtracking

- We need stronger algorithms with ideas from before: global order and message passing
- ABT (asynchronous backtracking)
- Agents are prioritized, messages pass from higher priority agents to lower priority agents
- Parallel execution
- Agents instantiate their variables concurrently and send their assigned values to the agents that are connected to them by outgoing links. All agents wait for and respond to messages. After each update of his assignment, an agent sends his new assignment along all outgoing links. An agent who receives an assignment (from the higher-priority agent of the link), tries to find an assignment for its variable that does not violate a constraint with the assignment it received

ABT communication

- Agent send messages ok?
- Agents stores received values into his data structure agent_view
- agent checks if his current assignment is consisted with his agent_view.
- If it is, the agent does nothing, otherwise it searchers for a new consistent value
- If the agent finds it, it assigns the found value to a variable and sends ok? message to all connected lower priority agents
- If the agent does not find it, it starts backtracking

ABT - backtracking

- The backtrack operation is executed by sending a Nogood message.
- Nogood is an inconsistent partial assignment (assignments of specific values to some of the variables that together violate the constraints on those variables)
- Nogood consists of A_i's agent_view
- The Nogood is sent to the agent with the lowest priority among the agents whose assignments are included in the inconsistent tuple in the Nogood.
- Agent A_i who sends a Nogood message to agent A_j assumes that A_j will change its assignment. Therefore, A_i removes from his agent_view the assignment of A_j and tries to find an assignment for A_j's variable that is consistent with the updated agent_view.

ABT properties

- Greedy hyper-resolution
- agents make tentative choices of a value for their variables, only generate Nogoods that incorporate values already generated by the agents above them in the order,
- communicates new values only to some agents and new Nogoods to only one agent.

ABT communication

when received (Ok?, (A_j, d_j)) do add (A_j, d_j) to agent_view check_agent_view

when received (Nogood, nogood) do add nogood to Nogood list forall $(A_k, d_k) \in$ nogood, if A_k is not a neighbor of A_i do add (A_k, d_k) to agent_view request A_k to add A_i as a neighbor check_agent_view

ABT check consistency

procedure check_agent_view

when agent_view and current_value are inconsistent do

if no value in D_i is consistent with agent_view then ↓ backtrack

else

select $d \in D_i$ consistent with *agent_view current_value* $\leftarrow d$ send (ok?, (A_i, d)) to lower-priority neighbors

ABT backtracking

procedure backtrack

 $nogood \leftarrow$ some inconsistent set, using hyper-resolution or similar procedure

if nogood is the empty set then

broadcast to other agents that there is no solution

terminate this algorithm

else

select $(A_j, d_j) \in nogood$ where A_j has the lowest priority in nogood send (Nogood, nogood) to A_j remove (A_j, d_j) from agent_view check_agent_view

ABT for c)

- Priority x₁, x₂, x₃
- They initially all start with a random value, e.g., all "blue"
- x₁ informs x₂ and x₃, x₂ informs x₃
 x₂ adds to its agent_view {x₁=blue}, x₃ adds {x₁=blue, x₂=blue}.
- * x_2 and x_3 have to check consistency with their own value
 - \times X₂ detects conflict, modifies its value to "red" and informs x₃
 - \times In that time x₃ detects conflict, modifies its value to "red", informs no one



Asynchronous backtracking for c)

- x₃ cannot find consistent value so using hyper-resolution it generates Nogood {x₁ = blue, x₂ = red}
- {red, blue} (c) $\neq \qquad X_1$ $X_1 \qquad \neq \qquad X_1$ $(x_2 \qquad \neq \qquad X_3$ {red, blue} $\neq \qquad {red, blue}$
- Sends this Nogood to x₂, because of lowest priority in Nogood
- Now x₂ cannot find consistent values and generates Nogood {x₁ = blue} and sends it to x₁.
- * x_1 detects inconsistency, modifies its value to "red" and informs x_2 and x_3
- As before, x₂ modifies its value to blue, x₃ cannot find consistent value and generates Nogood {x₁ = red, x₂ = blue},
- After that x₂ generates Nogood {x₁ = red} and sends it to x₁
- Now x₁ has Nogood {x₁ = blue} and {x₁ = red}, uses hyper-resolution to generate Nogood {}. Algorithm terminates by proclaiming that no solution exists.

Distributed Optimization

- agents shall, in a distributed fashion, optimize a global objective function
- We illustrate distributed path planning in directed graph with n nodes and m edges
- edge (a,b) has assigned a cost c(a,b);
- * objective: find a minimal cost path from the starting node s to any of the goal nodes $t \in T$.
- Applications: transport, telecommunications, planning
- Difference to standard algorithms (Dijkstra, Bellman-Ford) is a distributed approach (agents communicate only locally, each agent contributes to the globally optimal solution)

Asynchronous dynamic programming

- dynamic programming (incremental divide and conquer)
- if node x lies on a shortest path from s to t, then the portion of the path from s to x (and from x to t) must also be the shortest paths between s and x (x and t
- the shortest distance from any node *i* to the goal node *t* is represented with *h**(*i*).
- Shortest path from i to t via neighboring node j $f^{*}(i, j) = c(i, j) + h^{*}(j)$
- Shortest path from *i* via arbitrary neighboring node
 h*(i) = min_i f *(i, j)

Algorithm details

- Every node i stores a value h(i), as an approximation to h*(i)
- Initialization, each h(i)=∞,
- During execution, the h(i) values decrease and converge to the h*(i)
- Convergence takes one step for every node on the shortest path
- Weakness: we need an agent for every node

Pseudo code ADP for shortest path executed on every node

procedure ASYNCHDP (node i) if i is a goal node then $| h(i) \leftarrow 0$

else

initialize h(i) arbitrarily (e.g., to ∞ or 0)

repeat

forall neighbors j do $[f(j) \leftarrow w(i, j) + h(j)$ $h(i) \leftarrow \min_j f(j)$
ADP 1/4



ADP 2/4



ADP 3/4



ADP 4/4





- Algorithm LRTA* (learning real-time A*) uses one or more agents
- Heuristic search with improved heuristic
- Initialize: h(i)=0 (or any better informed admissible heuristics)
- Agent repeatedly executes an algorithm improving h(i)
- An example: a single agent execution

Pseudo code of LRTA*

procedure LRTA* $\iota \leftarrow s$ while i is not a goal node do foreach neighbor j do $f(j) \leftarrow w(i,j) + h(j)$ $i' \leftarrow \arg\min_i f(j)$ $\begin{array}{l} h(i) \leftarrow \max(\tilde{h}(i), f(i')) \\ i \leftarrow i' \end{array}$

An example:



One LRTA* agent 1/4



One LRTA* agent 2/4



One LRTA* agent 3/4



One LRTA* agent 4/4



LRTA*(2) (two agents) 1/3



LRTA*(2) (two agents) 2/3



LRTA*(2) (two agents) 3/3



Multiagent technologies

- Extensions to asynchronous backtracking
- Constraint satisfaction optimization
- Learning agents
- Game theory (cooperative and non-cooperative games

An exercise: simulate an execution of ADP and LRTA*

